PV168 Memory Model

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Retrospective

- Did you understand what was done in the last seminar?
- Was it intuitively understandable?
- Are you able to describe how it works?
- Did you find it difficult?

Monitor

- The goal is to prevent concurrent access using the critical section.
- Do not use multiple monitors on the same data.
- Each object instance and the class itself has a monitor.
 - Do not mix them.

Monitor

- Where to place synchronized?
 - $\circ\,$ in the method declaration
 - \circ in the static method declaration
 - $\circ\,$ in its own code block

Monitor

```
class Counter {
    private int value = 0;
    public synchronized int getValue() {
        return value++;
    public int getSameValue() {
        synchronized(this) {
            return value++;
```

Atomic numbers

- Useful for *just-a-counter* cases.
- Lower overhead
 - Monitors are quite expensive.
 - Can you say why?

Atomic numbers

```
class <u>Counter</u> {
    private AtomicInteger value = new <u>AtomicInteger(0);</u>
    public int getValue() {
        return value.getAndIncrement();
    }
}
```

But how monitors and atomics work?

Memory Model

- " In computing, a memory model describes the interactions of threads through memory and their shared use of the data.[1]
- Essential for any programming environment using parallelism.
- Consists of:
 - Language specification
 - Compiler specification (including optimizations)
 - Hardware (JVM in Java)

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Memory Model

- Specifies: [2]
 - $\circ\,$ synchronization points and actions
 - what is affected by synchronization points
 - $\circ\,$ what applies to memory access
 - before the synchronization point
 - after the synchronization point
- Defines *Sequential Consistency* and *Happens-before* relation.

Sequential Consistency

- " Compilers are allowed to reorder the instructions, when this does not affect the execution of the thread in isolation. [2]
- Also, the JIT may reorder, delay, or cache accesses.
- Also, the JVM may reorder, delay, or cache accesses.
- Also, the CPU may reorder, delay, or cache accesses.

"

Sequential Consistency

A = B = 0			
// thread 1 r2 = A; B = 1			
// thread 2 r1 = B; A = 2;			

What are the values of r1 and r2?
It may happen that r1 == 1 and r2 == 2.

Synchronization Points

- Java has 4 types:
 - volatile
 - Atomic classes
 - o synchronized
 - $\circ\,$ starts and joins of threads

Actions

- Action types:
 - \circ Load
 - Store
 - Synchronization
- Loads without any store are trouble-free.
- Any store requires a proper synchronization.
 - Including all load accesses!
- Synchronization locks and unlocks the monitor.

Shared variables

- Shared variables (e.g. affected by the *Memory Model*):
 - \circ instance fields
 - static fields
 - array elements
- Variables unaffected by the *Memory Model*:
 - \circ local variables
 - methods arguments, catched exceptions

Happens-before

- Transitive relationship of actions.
 - " If one action happens-before another, then the first is visible to and ordered before the second [3]
- Applies to all synchronization points.
- Partially applies to all final fields.
 - After construction, no further synchronization is needed.

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volatile

- Implements the memory barrier:
 - accesses to volatile s cannot be reordered
 - subsequent accesses cannot be reordered before volatile load
 - prior accesses cannot be reordered after volatile store
- Numeric operations are not atomic (++, --)
 - Store of 64 bit values is done in two steps.

• DO NOT USE IT UNLESS YOU KNOW WHAT YOU ARE DOING!

volatile

```
class VolatileExample {
```

```
int x = 0;
volatile boolean flag = false;
public void writer() {
    x = 42;
    flag = true;
}
```

```
public void reader() {
    if (flag == true) {
        System.out.println(x); // guaranteed to see 42.
```

Atomic classes

- Useful for numeric counters, gauges, and indicators.
- Thread safe for all basic numeric operations.
 - Increments, decrements, assignment
- "Synchronizes properly as you would expect."
 - Almost the same rules applies as for volatile s.
 - Allows finer specification of memory ordering.

• DO NOT USE FOR OTHER THINGS THAN COUNTERS UNLESS YOU KNOW WHAT YOU ARE DOING!

Atomic classes

```
class AtomicExample {
    private int value = 0;
    private AtomicBoolean flag = new AtomicBoolean(false);
    public void writer() {
        value = 42;
        flag.set(true);
    }
```

```
public void reader() {
```

while (!flag.get()) {} // busy wait here
System.out.println(value); // guaranteed to see 42.

synchronized

- When used in a method signature, it uses the monitor of the instance.
 - Or monitor of the class in case the method is static.
- When used in the body of a method, it uses the monitor of the specified object.

public synchronized void doSomething() { /* ... */ }
public void doSomethingElse() {
 synchronized(this.attribute) {/* ... */}

Threads

- Thread start and join have similar memory semantics as synchronized :
 - Any store which *happened-before* the start of the thread is visible by the thread.
 - Any store which *happened-before* the join of the *about-to-join* thread is visible in the joinee thread.

Threads

```
class ThreadExample {
   public void run() {
      Integer value = 42;
      var t = new Thread(() -> value *= 2 );
      // value is 42
      t.start();
      t.join();
      // value is guaranteed to be 84;
}
```



- Special rules apply to final fields: [4]
 - Set the final fields for an object in that object's constructor.
 - Do not write a reference to the object being constructed in a place where another thread can see it before the object's constructor is finished.
 - If this is followed, then when the object is seen by another thread, that thread will always see the correctly constructed version of that object's final fields.

Double-Checked Locking

Is this correct?

Double-Checked Locking Is Broken And there is no way how it can be fixed.

How To Test The Correct Usage?

- Unit tests are insufficient.
 - In fact, no commonly used automatization approaches can validate the correctness.
- To prove the correctness, one must provide a formal proof.
 - This is the reason airplanes work in strictly single-threaded environments.

Without correct synchronization, very strange, confusing and counterintuitive behaviors are possible.